
1 More Running Time

Give the worst case and best case running time in $\Theta(\cdot)$ notation in terms of M and N .

(a) Assume that `comeon()` $\in \Theta(1)$ and returns a boolean.

```
1 for (int i = 0; i < N; i += 1) {
2     for (int j = 1; j <= M; ) {
3         if (comeon()) j += 1;
4         else          j *= 2;
5     }
6 }
```

2 Recursive Running Time

For the following recursive functions, give the worst case and best case running time in the appropriate $O(\cdot)$, $\Omega(\cdot)$, or $\Theta(\cdot)$ notation.

(a) Give the running time in terms of N .

```
1 public void andslam(int N) {
2     if (N > 0) {
3         for (int i = 0; i < N; i += 1) {
4             System.out.println("datboi.jpg");
5         }
6         andslam(N / 2);
7     }
8 }
```

(b) Give the running time for `andwelcome(arr, 0, N)` where N is the length of the input array `arr`.

```
1 public static void andwelcome(int[] arr, int low, int high) {
2     System.out.print("[ ");
3     for (int i = low; i < high; i += 1) {
4         System.out.print("loyal ");
5     }
6     System.out.println("]");
7     if (high - low > 0) {
8         double coin = Math.random();
9         if (coin > 0.5) {
10            andwelcome(arr, low, low + (high - low) / 2);
11        } else {
12            andwelcome(arr, low, low + (high - low) / 2);
13            andwelcome(arr, low + (high - low) / 2, high);
14        }
15    }
16 }
```

(c) Give the running time in terms of N .

```
1 public int tothe(int N) {
2     if (N <= 1) {
3         return N;
4     }
5     return tothe(N - 1) + tothe(N - 1);
6 }
```

(d) *Extra Hard!* Give the running time in terms of N

```
1 public static void spacejam(int N) {
2     if (N == 1) {
3         return;
4     }
5     for (int i = 0; i < N; i += 1) {
6         spacejam(N - 1);
7     }
8 }
```

3 Hey you watchu gon do?

For each example below, there are two algorithms solving the same problem. Given the asymptotic runtimes for each, is one of the algorithms **guaranteed** to be faster? If so, which? And if neither is always faster, explain why. Assume the algorithms have very large input (so N is very large).

- (a) Algorithm 1: $\Theta(N)$, Algorithm 2: $\Theta(N^2)$
- (b) Algorithm 1: $\Omega(N)$, Algorithm 2: $\Omega(N^2)$
- (c) Algorithm 1: $O(N)$, Algorithm 2: $O(N^2)$
- (d) Algorithm 1: $\Theta(N^2)$, Algorithm 2: $O(\log N)$
- (e) Algorithm 1: $O(N \log N)$, Algorithm 2: $\Omega(N \log N)$

Would your answers above change if we did not assume that N was very large?

4 More Extra Problems [Final FA15]

If you have time try, to answer this challenge question. For each answer true or false. If true, explain why and if false provide a counterexample.

(a) If $f(n) \in O(n^2)$ and $g(n) \in O(n)$ are positive-valued functions (that is for all n , $f(n), g(n) > 0$), then $\frac{f(n)}{g(n)} \in O(n)$.

(b) If $f(n) \in \Theta(n^2)$ and $g(n) \in \Theta(n)$ are positive-valued functions, then $\frac{f(n)}{g(n)} \in \Theta(n)$.